

Atiyah described mathematics as the “science of analogy.” In this vein, the purview of category theory is *mathematical analogy*. Category theory provides a cross-disciplinary language for mathematics designed to delineate general phenomena, which enables the transfer of ideas from one area of study to another. The category-theoretic perspective can function as a simplifying abstraction, isolating propositions that hold for formal reasons from those whose proofs require techniques particular to a given mathematical discipline. [1]

**Definition 0.1.** A **category** is made of the following **data**

- a collection of **objects**:  $a, b, c, \dots$
- a collection of **morphisms**:  $f, g, h, \dots$

<sup>1</sup> such that

- each morphism has a **domain** and a **codomain**:  $f : a \rightarrow b$
- each object has a designated **identity morphism**:  $1_a : a \rightarrow a$
- for any pair of morphisms  $f, g$  such that  $\text{cod}f = \text{dom}g$  there exists a designated **composite morphism**  $g \circ f$

$$\begin{array}{ccccc}
 & & g \circ f & & \\
 & \frown & & \searrow & \\
 a & \xrightarrow{f} & b & \xrightarrow{g} & c
 \end{array}$$

this **data** obeys the following axioms

- for any  $f : a \rightarrow b$

$$f \circ 1_a = f = 1_b \circ f$$

- for any composable morphisms  $f, g, h$

$$(h \circ g) \circ f = h \circ (g \circ f) = h \circ g \circ f$$

---

<sup>1</sup>here collection is used to avoid Russel's paradox

...the whole concept of a category is essentially an auxiliary one; our basic concepts are essentially those of a **functor** and of a **natural transformation**... [2]

See: [2], [1] Ch.1.1, [3] Ch.1.1, [4]. Here are some examples.

---

**Definition 0.2.** A category is **small** if the collection of its morphisms is a set (it follows that its objects form a set too).

**Definition 0.3.** A category is **locally small** if the collection of morphisms between any pair of fixed objects is a set.

**Definition 0.4.** In a locally small category  $\mathcal{C}$ , for any pair of objects  $a, b \in \mathcal{C}$  the **hom-set**  $\mathcal{C}(a, b)$  is the set of morphisms  $a \rightarrow b$ .

---

**Definition 0.5.** A morphism  $f : a \rightarrow b$  is an **isomorphism** iff there exists  $g : b \rightarrow a$  such that  $f \circ g = 1_b$  and  $g \circ f = 1_a$ . Then the objects  $a, b$  are **isomorphic**  $a \cong b$ .

In a category there is no notion of equality of objects (they can be isomorphic at most), there is equality of morphisms though (which is used extensively in commuting diagrams).

For a discussion on equality in mathematics see [5].

It is easy to see that if  $a \cong b$  then if there exists a morphism  $c \rightarrow a$  (or  $a \rightarrow c$ ) then there must exist  $c \rightarrow b$  (or  $b \rightarrow c$ ). Thus for two objects being isomorphic means having the same relationship with the rest of the category.

---

[1] E. Riehl, *Category Theory in Context* (Dover, 2015).

[2] S. M. Lane, *Categories for the Working Mathematician* (Springer New York, 1971).

[3] P. Perrone, *Notes on Category Theory with Examples from Basic Mathematics*, (2019).

[4] B. Milewski, *Category Theory for Programmers* (2019).

[5] B. Mazur, *When Is One Thing Equal to Some Other Thing?*, (2006).

---